Semantic conflict detection via dynamic analysis

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Centro de Informática

The software development process is strongly **collaborative**











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Semantic conflict (behavioral)

When the merge of contributions from 2 developers results in unplanned behavior when compared to the behavior of the versions that originated it



Unplanned behavior ~ Interference



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Overriding Assignment

occurs when additions or modifications to one of the merge branches involve writing to a state element that is also associated with a write operation due to the contribution of the other branch





Overriding Assignment

occurs when additions or modifications to one of the merge branches **involve** writing to a **state element** that is also associated with a write operation due to the contribution of the other branch



[...] **involve** writing to a state element

1	<pre>class Text {</pre>
2	constructor() {
3	<pre>this.text = "";</pre>
4	<pre>this.fixes = 0;</pre>
5	<pre>this.comments = 0;</pre>
6	}
7	generateReport() {
8	<pre>countDuplicatedWhitespaces(); // Left</pre>
9	<pre>countComments();</pre>
10	<pre>countDuplicatedWords(); // Right</pre>
11	}



[...] involve writing to a state element







Semantic conflict detection via dynamic analysis

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Abstract

During collaborative software development, a semantic conflict may occur when the individual behavior expected by different developers is no longer preserved after merging their branches. While potential semantic conflicts are not captured via textual merge tools, different approaches have already been proposed based on static analysis or automated test generation to verify behavioral changes given a merge scenario. However, these approaches share some limitations regarding scalability and reporting false positives and negatives. Trying to address these limitations, in this work, we assess the detection of conflicts by focusing on overriding textual conflicts, behavioral semantic conflicts represent the behavioral differences between a program's pre- and postmerge versions, impacting how the software works when executed (either during test suites or even in production environments). The motivation to detect this kind of conflict is due to the divergence of behavior that occurs even when the integration is textually and syntactically valid [7], resulting in a scenario that incurs costs for software development when not noticed before execution.

The expected behavior intended by a software contributor is hard to assess and, in this context, can be approximated into program specifications about how it should work, while its unexpected change (potential semantic conflict) can be





Why and how?

Semantic conflict detection via dynamic analysis

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Abstract

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Related work

Static Analysis

Conservative; Potential drop in accuracy

Unit Test Generation

Run on each version of merge; potential increase of false negatives



Related work

Static Analysis

Conservative; Potential drop in accuracy

Unit Test Generation

Run on each version of merge; potential increase of false negatives

Semantic conflict detection with overriding assignment analysis

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levar à introdução de bugs no código, influenciando negativamente na qualidade do produto final. Horwitz et al. [18] especificaram for-

malmente os conflitos semânticos: duas contribuições advindas de

RESUMO

Developers typically work collaboratively and often need to embed their code into a major version of the system. This process can cause

Detecting Semantic Conflicts using Static Analysis

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ABSTRACT

Version control system tools empower developers to independently work on their development tasks. These tools also facilitate the integration of changes through merging operations, and report textual conflicts. However, when developers integrate their changes, they might encounter other types of conflicts that are not detected by current merge tools. In this paper, we focus on dynamic semantic conflicts, which occur when merging reports no textual conflicts but results in undesired interference- causing unexpected program behavior at runtime. To address this issue, we propose a technique that explores the use of static analysis to detect interference when merging contributions from two developers. We evaluate our technique using a dataset of 99 experimental units extracted from merge scenarios. The results provide evidence that our technique presents significant interference detection capability. It outperforms, in terms of F1 score and recall, previous methods at roly on dynamic analysis for detecting semantic conflicts but

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Even worse, textual merge tools also can't detect dynamic semantic confile (6, 8, 12, 19, 24, 27, 33, 40, 41), like when the changes made by one developer affect a state element that is accessed by code changed by another developer, who assumed a state inviraint that no longer holds after merging. In such cases, textual integration is automatically performed generating a merged point state and the success for this program, but its execution reveals unexpected behavior.¹ Following Horwitz et al. [19], we put this more formably by considering dynamic semantic conflicts to be *undesired interference*, where interference is defined as follows: separate changes L and R to abser program B interfere when the integrated changes does not preserve the changed behavior of L or R, or the unchanged behavior of B.

As dynamic semantic conflicts, hereafter simply semantic conflicts, can negatively impact development productivity and the quality of software products, researchers [12, 19, 33] have proposed techniques to detect them. In fact, as developer's desire (specific es e não sabemos a intenção real dos entes já foram utilizadas, como, por 8]. Porém, essas tendem a apresentar titvos, pois verificam se há interferênicos na execução do código testado. se algoritmos de análise estática para so [5, 14, 18]. mas esses são baseados

um programa Base, originam um conflito ões que as versões se propõem a cumfeitas na versão integrada Merge.¹ Na a existência dos conflitos semânticos,



Related work

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Unit Test Generation

Run on each version of merge; potential increase of false negatives

Detecting Semantic Conflicts via Automated Behavior Change Detection

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Abstract-Branching and merging are common practices in collaborative software development. They increase developer productivity by fostering teamwork, allowing developers to independently contribute to a software project. Despite such benefits, branching and merging comes at a cost-the need to merge software and to resolve merge conflicts, which often occur in practice. While modern merge techniques, such as 3-way or structured merge, can resolve many such conflicts automatically. they fail when the conflict arises not at the syntactic, but the semantic level. Detecting such conflicts requires understanding the behavior of the software, which is beyond the capabilities of most existing merge tools. As such, semantic conflicts can only be identified and fixed with significant effort and knowledge of the changes to be merged. While semantic merge tools have been proposed, they are usually heavyweight, based on static analysis, and need explicit specifications of program behavior. In this work, we take a different route and explore the automated creation of unit tests of partial specifications to detect unwanted behavior changes (conflicts) when merging software.

We systematically explore the detection of semantic conflicts through unit-test generation. Relying on a ground-truth dataset of 38 software merge scenarios, which we extracted from GitHub, we manually analyzed them and investigated whether

changes to be merged. This can negatively affect development productivity, and even compromise software quality in case developers incorrectly fix conflicts [17], [6], [18]. To avoid dealing with merge conflicts, developers sometimes even adopt risky practices, such as rushing to finish changes first [19], [17] and partial check-ins [20]. Similarly, partially motivated by the need to reduce merge conflicts, development teams have been adopting techniques such as trunk-based development [21], [22], [23] and feature togets [24], [21], [25], [26].

Although these practices might reduce the occurrence of merge conflicts, there is no evidence that they are effective for resolving or even detecting so-called test [8] and production conflicts, which are only observed when running project tests and using the system in production. As such, they are more serious, because they reveal software failures. In fact, some of the practices mentioned above might even aggravate the costs of test and production conflicts, which are special kinds of what we here call semantic conflicts. To make matters worse, we expect semantic conflicts to cost more than merge conflicts, as they are often bardet to detect and resolve, and might end



Available resources

Dynamic Analysis

Occurs at runtime

- Does not demand assertions about the program
- 🔻 It is only executed on the post-merge code

README	ৰ্ক্য Apache-2.0 license	
Jalan	gi2	

Introduction

README

Jalangi2 is a framework for writing dynamic analyses for JavaScript. Jalangi1 is still available at https://github.com/SRA-SiliconValley/jalangi, but we no lenger plan to levelop it. Jalangi2 does not support the record/replay feature of Jalangi1. In the Jalangi2 distribution you will find several analyses:

- an analysis to track NaNs.
- · an analysis to check if an undefined is concatenated to a string.
- · Memory analysis: a memory-profiler for JavaScript and HTML5.
- · DLint: a dynamic checker for JavaScript bad coding practices.
- JITProf: a dynamic JIT-unfriendly code snippet detection tool.
- · analysisCallbackTemplate.js: a template for writing a dynamic analysis.
- and more ...

See our tutorial slides for a detailed overview of Jalangi and some client analyses.

Requirements

We have tested Jalangi on Mac OS X with Chromium browser. Jalangi should work on Mac OS 10.7, Ubuntu 11.0 and higher and Windows 7 or higher. Jalangi will NOT work with IE.

- · Node.js available at https://nodejs.org/en/download/releases/. We primarily test Jalangi with the Node LTS version (currently v12).
- · Chrome browser if you need to test web apps.
- Python (http://python.org) version 3.x.

On Windows you need the following extra dependencies:

Install Microsoft Visual Studio 2010 (Free oversion version is fine)



Jalangi2 Analysis



Overriding Assignment Analysis



Custom OA Analysis











Algorithm 1: Overriding Assignment Detection Data: A script s Result: A list of overriding assignment interferences 1 for event \in s do if $isWrite(event) \lor isPutFieldPre(event)$ then 2 **if** *isFromSomeBranch*(*event*) ∨ 3 \neg isFunctionCallStackEmpty() then addAssignmentToCurrentBranch(event) 4 if hasAssignmentFromOtherBranch(event) then 5 updateInterferences(event) 6 removeAssignmentFromOtherBranch(event) 7 else 8 removeAssignmentFromAllBranches(event) 9 if isInvokeFunPre(event) then 10 **if** *isFromSomeBranch(event)* ∨ 11 ¬isFunctionCallStackEmpty() then if isArrayInplaceMethod(event) then 12 handleAssignedIndices(event) 13 else 14 functionCallStack.push(event) 15 if isInvokeFun(event) then 16 functionCallStack.pop() 17 if *isEndExecution(event)* then 18 logResults() 19





	0 0 0
	Data: A script s
	Result: A list of overriding assignment interferences
1	for $event \in s$ do
2	if isWrite(event) ∨ isPutFieldPre(event) then
3	if $isFromSomeBranch(event) \lor$
	¬isFunctionCallStackEmpty() then
4	addAssignmentToCurrentBranch(event)
5	if hasAssignmentFromOtherBranch(event) then
6	updateInterferences(event)
7	removeAssignmentFromOtherBranch(event)
8	else
9	removeAssignmentFromAllBranches(event)
10	if isInvokeFunPre(event) then
11	if $isFromSomeBranch(event) \lor$
	¬isFunctionCallStackEmpty() then
12	if isArrayInplaceMethod(event) then
13	handleAssignedIndices(event)
14	else
15	functionCallStack.push(event)
16	if isInvokeFun(event) then
17	functionCallStack.pop()
18	if isEndExecution(event) then
19	logResults()





	Data: A script s		
	Result: A list of overriding assignment interferences		
1	for $event \in s$ do		
2	if <i>isWrite</i> (<i>event</i>) ∨ <i>isPutFieldPre</i> (<i>event</i>) then		
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	\neg isFunctionCallStackEmpty() then		
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Function call stack



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	Result: A list of overriding assignment interferences			
1	for event \in s do			
2	if $isWrite(event) \lor isPutFieldPre(event)$ then			
3	if <i>isFromSomeBranch(event)</i> ∨			
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	Date: A seriet a		
	Data: A script s		
	Result: A list of overriding assignment interferences		
1	$ \begin{array}{c} \text{for even} \in S \text{ do} \\ \text{for even} \\ \text{for even} \in S \text{ do} \\ \text{for even} \\ \text{for even}$		
2	if isWrite(event) V isPutFieldPre(event) then		
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Interferences Left assignments b

Function call stack



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Interferences

Left assignments

b

Right assignments

Function call stack



	Data: A script s
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1	for $event \in s$ do
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Function call stack



_	10010		
	Data: A script s		
	Result	: A list of overriding assignment interferences	
1	for eve	$ent \in s do$	
2	if	$isWrite(event) \lor isPutFieldPre(event)$ then	
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16	if	isInvokeFun(event) then	
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func1

	Data:	A script s							
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19	logResults()							





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Algorithm 1: Overriding Assignment Detection Data: A script s Result: A list of overriding assignment interferences 1 for event \in s do if $isWrite(event) \lor isPutFieldPre(event)$ then 2 **if** *isFromSomeBranch(event)* ∨ 3 \neg isFunctionCallStackEmpty() then addAssignmentToCurrentBranch(event) 4 if hasAssignmentFromOtherBranch(event) then 5 updateInterferences(event) 6 *removeAssignmentFromOtherBranch(event)* 7 else 8 removeAssignmentFromAllBranches(event) 9 if isInvokeFunPre(event) then 10 **if** *isFromSomeBranch(event)* ∨ 11 ¬isFunctionCallStackEmpty() then if isArrayInplaceMethod(event) then 12 handleAssignedIndices(event) 13 else 14 functionCallStack.push(event) 15 if isInvokeFun(event) then 16 functionCallStack.pop() 17 if *isEndExecution(event*) then 18 logResults() 19





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_							
	Data: A script s						
	Result: A list of overriding assignment interferences						
1	for event \in s do						
2	if $isWrite(event) \lor isPutFieldPre(event)$ then						
3	if <i>isFromSomeBranch</i> (<i>event</i>) ∨						
	\neg isFunctionCallStackEmpty() then						
4	addAssignmentToCurrentBranch(event)						
5	if hasAssignmentFromOtherBranch(event) then						
6	updateInterferences(event)						
7	removeAssignmentFromOtherBranch(event)						
8	else						
9	removeAssignmentFromAllBranches(event)						
10	if isInvokeFunPre(event) then						
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```
"type": "OVERRIDING_ASSIGNMENT_CONFLICT",
```

"label": "Overriding Assignment Conflict",

"body": {

ι

"description": "Interference detected on 3_b:\nBranch L at line 54 (/index.js:54:5:54:6)\nBranch R at line 50 (/index.js:50:9:50:10) with interprocedural stack: [{\"id\":1241,\"name\":\"func1\",\"location\":\"

```
(/index.js:56:1:56:8)\",\"branch\":\"R\",\"beforeInvoke\":true}]",
    "interference": {
     "previousAssignment": {
       "id": 3,
      "name": "b",
      "location": "(/index.js:54:5:54:6)",
      "branch": "L",
       "isObject": false
     },
     "currentAssignment": {
       "id": 3,
      "name": "b",
      "location": "(/index.js:50:9:50:10)",
       "branch": "R",
       "isObject": false,
      "functionCallStack": [
         "id": 1241,
         "name": "func1",
         "location": "(/index.js:56:1:56:8)",
         "branch": "R",
         "beforeInvoke": true
```







Evaluation & Results













... IfBranchConflictSample /* Generated from Java with JSweet 3.0.0 - http://www.jsweet.org */ var br; (function (br) { var unb: (function (unb) { var cic: (function (cic) { var analysis: (function (analysis) { var samples: (function (samples) { var ioa: (function (ioa) { var IfBranchConflictSample = /** @class */ (function () { function IfBranchConflictSample() { if (this.x === undefined) { this.x = 0: l IfBranchConflictSample.main = function () { var m = new IfBranchConflictSample(); m.x = 0;m.foo(); }; /*private*/ IfBranchConflictSample.prototype.foo = function () { if $(this.x \ge 0)$ { this.x = 1;3 else { var a = 0;}: return IfBranchConflictSample; }()); ioa.IfBranchConflictSample = IfBranchConflictSample; IfBranchConflictSample["__class"] = "br.unb.cic.analysis.samples.ioa.IfBranchConflictSample"; })(ioa = samples.ioa || (samples.ioa = {})); })(samples = analysis.samples || (analysis.samples = {})); })(analysis = cic.analysis || (cic.analysis = {})); })(cic = unb.cic || (unb.cic = {})); })(unb = br.unb || (br.unb = {})); })(br || (br = {})); br.unb.cic.analysis.samples.ioa.IfBranchConflictSample.main():

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```
IfBranchConflictSample
package br.unb.cic.analysis.samples.ioa;
// Conflict: [left, main():10] --> [right, foo():16]
public class IfBranchConflictSample {
   private int x;
   public static void main() {
       IfBranchConflictSample m = new IfBranchConflictSample();
       m.x = 0; // LEFT
       m.foo(); // RIGHT
    }
   private void foo() {
       if (x >= 0) {
            x = 1;
       } else {
           int a = 0;
           // System.out.println(x)
        }
    }
}
```

```
}
IfBranchConflictSample.main = function () {
    var m = new IfBranchConflictSample();
    m.x = 0;
    m.foo();
};
/*private*/ IfBranchConflictSample.prototype.foo = function () {
    if (this.x >= 0) {
        this.x = 1;
    }
    else {
        var a = 0;
    }
};
//(Cic = unb.cic || (unb.cic = {}));
```

```
})(unb = br.unb || (br.unb = {}));
})(br || (br = {}));
br.unb.cic.analysis.samples.ioa.IfBranchConflictSample.main();
```



Test Results

In translated validation scenarios

Transpilation Result	Amount of tests	Success Rate
TRANSPILATION_DIVERGENCE	2	-
CONCEPTUALLY_ADAPTED	14	<mark>100%</mark>
SUCCESS	<mark>55</mark>	<mark>100%</mark>













In open source merge scenarios



























With an additional overriding assignment analysis, this time disregarding branches or whether it was modified,

83% of **159** scenarios showed no override of any state element







With an additional overriding assignment analysis, this time disregarding branches or whether it was modified,

83% of 159 scenarios showed no override of any state element 83% true

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$\bullet \bullet \bullet$

Interference detected on 5_inject:

Branch L at (.../Nightmare/lib/actions.js:100:1:120:2)
Branch R at (.../Nightmare/lib/actions.js:226:1:240:2)

```
exports.inject = function(type, file, done){
       debug('.inject().ing a file');
       var startTag, endTag;
       if ( type !== "js" && type !== "css" ){
          debug('unsupported file type in .inject()');
         done();
       if (type === "js"){
          startTag = "<script>";
         endTag = "</script>";
        } else if (type === "css"){
         startTag = "<style>";
         endTag = "</style>":
       var self = this;
       this.page.getContent( function(pageContent){
          var injectedContents = fs.readFileSync(file);
          self.page.setContent(pageContent + startTag + inject
       });
120
```

100% de positivos verdadeiros

Apenas 1 positivo, que é verdadeiro

```
@param {String} type
 * @param {String} file
 * @param {Function} done
exports.inject = function(type, file, done){
 var startTag, endTag;
 if (type === "js"){
   startTag = "<script>";
   endTag = "</script>";
 } else {
   startTag = "<style>";
   endTag = "</style>";
 var self = this;
 this.page.getContent( function(pageContent){
   var injectedContents = fs.readFileSync(file);
   self.page.setContent(pageContent + startTag + inject
 });
```









Overhead analysis

The native Jalangi2 instrumentation and analyses represent the major overhead









In the future...

More algorithms to detect conflicts beyond OA

1	var	ar	r =	[0,	0,	0,	0,	0]	
2	arr	=	[1,	1,	1,	1,	1,	1,	1]
3	cons	ol	e.lc	g('	hel	10')		
4	arr	[1]	=	1					
5	<pre>console.log('world!')</pre>								
6	arr	[2]	= :	2					
7	cons	ol	e.lc	g (a	rr)				



In the future...

More algorithms to detect conflicts beyond OA

1	var	ar	r =	[0,	0,	0,	0,	0]	
2	arr	=	[1,	1,	1,	1,	1,	1,	1]
3	cons	o lo	e.lo	og (′	hel	lo ′)		
4	arr	[1]	=	1					
5	<pre>console.log('world!')</pre>								
6	arr	[2]	= :	2					
7	cons	sol	e.lo	g(a	rr)				

Plug the solution into code integration pipelines





```
function aumentarContador() {
```

```
...
}
function iniciarProcesso() {
    ...
}
module.exports = {
    aumentarContador, iniciarProcesso
}
```

Expand analysis action on scripts instructions

Fuzzers, generating tests without assertions, or reusing parameters present in tests from the repositories



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Thank you!





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